# Block Factor-width-two Matrices and Their Applications to Semidefinite and Sum-of-squares Optimization 

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## Outline

Introduction: inner/outer approximations for SDPs

A new class of block factor-width-two matrices

Applications to SDPs and SOS optimization

Conclusions

## Introduction

A primal standard SDP is in the form of

$$
\begin{array}{ll}
p^{\star}=\min _{X} & \langle C, X\rangle \\
\text { subject to } & \left\langle A_{i}, X\right\rangle=b_{i}, i=1, \ldots, m \\
& X \in \mathbb{S}_{+}^{n}
\end{array}
$$

SDPs are more powerful than LP or SOCP

- Applications: control theory, polynomial optimization, machine learning, power grid, robotics, etc.



## Introduction

A primal standard SDP is in the form of

$$
\begin{aligned}
p^{\star}=\min _{X} & \langle C, X\rangle \\
\text { subject to } & \left\langle A_{i}, X\right\rangle=b_{i}, i=1, \ldots, m, \\
& X \in \mathbb{S}_{+}^{n} .
\end{aligned}
$$

## However, SDPs are much more expensive to solve than LP or SOCP

- Standard interior-point methods scale as $\mathcal{O}\left(m n^{3}+m^{2} n^{2}\right)$ per iteration
- LPs with millions of variables and constraints can be solved reliably.
- General-purpose solvers cannot efficiently handle large SDP problems ( $n \approx 1000$, and $m$ : a few thousands)
- Exploiting sparsity and structures for improving efficiency is an active research topic ${ }^{1,2}$.

[^0]
## Something simpler: inner/outer approximations

## Inner approximations

- Suppose we have a simpler cone $\mathcal{K} \subset \mathbb{S}_{+}^{n}$. Solving an instance of

$$
\begin{aligned}
\min _{X} & \langle C, X\rangle \\
\text { subject to } & \left\langle A_{i}, X\right\rangle=b_{i}, i=1, \ldots, m \\
& X \in \mathcal{K} \subset \mathbb{S}_{+}^{n}
\end{aligned}
$$

gives us an upper bound on $p^{\star}$.

## Outer approximations

- Suppose we have a simpler cone $\mathbb{S}_{+}^{n} \subset \hat{\mathcal{K}}$. Solving an instance of

$$
\begin{aligned}
\min _{X} & \langle C, X\rangle \\
\text { subject to } & \left\langle A_{i}, X\right\rangle=b_{i}, i=1, \ldots, m \\
& X \in \hat{\mathcal{K}} .
\end{aligned}
$$

gives us a lower bound on $p^{\star}$.

## Which cones to use?

Ahmadi and Majumdar ${ }^{3}$ considered the cones of diagonally dominant and scaled diagonally dominant matrices

- A symmetric matrix $A \in \mathbb{S}^{n}$ is diagonally dominant if

$$
a_{i i} \geq \sum_{j \neq i}\left|a_{i j}\right|, \quad i=1, \ldots, n
$$

- A symmetric matrix $A \in \mathbb{S}^{n}$ is scaled-diagonally dominant if there exists a diagonal matrix $D$ with nonnegative entries, such that
$D A D$ is diagonally dominant
We denote

$$
\begin{aligned}
\mathcal{D D}_{n} & =\left\{X \in \mathbb{S}^{n} \mid X \text { is diagonally dominant }\right\} \subset \mathbb{S}_{+}^{n} \\
\mathcal{S D D}{ }_{n} & =\left\{X \in \mathbb{S}^{n} \mid X \text { is scaled diagonally dominant }\right\} \subset \mathbb{S}_{+}^{n}
\end{aligned}
$$

- Linear optimization over $\mathcal{D} \mathcal{D}_{n}$ is an LP;
- Linear optimization over $\mathcal{S D D}_{n}$ is an SOCP;

[^1]
## Approximation and Conservativeness

$$
\begin{aligned}
\min _{X} & \langle C, X\rangle \\
\text { subject to } & \left\langle A_{i}, X\right\rangle=b_{i}, i=1, \ldots, m, \\
& X \in \mathcal{S D D}_{n}\left(\text { or } \mathcal{D D}_{n}\right)
\end{aligned}
$$

- First: the resulting upper bound might be very conservative.


Figure: Boundary of $(x, y)$ such that $I_{6}+x A+x B$ is PSD, SDD, or DD.

- Second: using $\mathcal{S D D}{ }_{n}$ requires a number of $\mathcal{O}\left(n^{2}\right)$ small SOCP constraints, which might be an issue for large $n$.


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## Factor-width-two matrices

## Lemma

$\mathcal{S D D}_{n}$ is equivalent to the cone of factor-width-two matrices (Boman, Erik G., et al. 2005)

$$
\mathcal{S D D}_{n}=\left\{\sum_{i \in I} x_{i} x_{i}^{\top} \mid x_{i} \in \mathbb{R}^{n}, \operatorname{supp}\left(x_{i}\right) \leq 2\right\} .
$$

- Denote the cone of $n \times n$ factor-width-two matrices as $\mathcal{F} \mathcal{W}_{2}^{n}$ :

$$
\mathcal{S D D}_{n}=\mathcal{F} \mathcal{W}_{2}^{n} .
$$

- Another interpretation: $P \in \mathcal{F} \mathcal{W}_{2}^{n}$ if and only if there exists $X_{i j} \in \mathbb{S}_{+}^{2}$ s.t.

$$
P=\sum_{1 \leq i<j \leq n} E_{i j}^{\top} X_{i j} E_{i j},
$$

where $E_{i j}=\left[\begin{array}{c}E_{i} \\ E_{j}\end{array}\right] \in \mathbb{R}^{2 \times n}, i \neq j$, and $E_{i}$ is a row basis vector with 1 at the $i$-th entry.

$$
E_{i}=\left[\begin{array}{lllll}
0 & \ldots & 1 & \ldots & 0
\end{array}\right] \in \mathbb{R}^{1 \times n} .
$$

## Block factor-width-two matrices

A linear optimization problem over the $F W_{2}^{n}$ cone can be written as an SDP over the cone product

$$
\underbrace{\mathbb{S}_{+}^{2} \times \ldots \times \mathbb{S}_{+}^{2}}_{\binom{n}{2}}
$$



Figure: Illustration of (block) factor-width-two decomposition

- For $A \in \mathcal{S D D} \mathcal{D}_{n}$ or $\mathcal{F W}_{2}^{n}$, each black square represents a scalar $a_{i j} \in \mathbb{R}$.
- Key idea - block extension: how about each black square represents a submatrix?

$$
A_{i j} \in \mathbb{R}^{k_{i} \times k_{j}}
$$

## Block-partitioned matrices

Given a matrix $A \in \mathbb{R}^{n \times n}$, we say a set of integers $\alpha=\left\{k_{1}, k_{2}, \ldots, k_{p}\right\}$ with $k_{i} \in \mathbb{N}(i=1, \ldots, p)$ is a partition of $A$ if $\sum_{i=1}^{p} k_{i}=n$, and $A$ is partitioned as

$$
\left[\begin{array}{ccc}
A_{11} & A_{12} & \ldots A_{1 p} \\
A_{21} & A_{22} & \ldots A_{2 p} \\
\vdots & \vdots & \ddots \\
A_{p 1} & A_{p 2} & \ldots A_{p p}
\end{array}\right]
$$

with $A_{i j} \in \mathbb{R}^{k_{i} \times k_{j}}, \forall i, j=1, \ldots, p$.

(a)

(b)

(c)

Figure: Different partitions for a $6 \times 6$ matrix: (a) $\alpha=\{1,1,1,1,1,1\}$, (b) $\beta=\{2,2,2\}$, (c) $\gamma=\{4,2\}$. From right to left, we get coarser partitions, i.e. $\alpha \sqsubseteq \beta \sqsubseteq \gamma$.

## Block factor-width-two matrices

## Definition

A symmetric matrix $Z \in \mathbb{S}^{n}$ with partition $\alpha=\left\{k_{1}, k_{2}, \ldots, k_{p}\right\}$ belongs to the class of block factor-width-two matrices, denoted as $\mathcal{F} \mathcal{W}_{\alpha, 2}^{n}$, if and only if

$$
\begin{equation*}
Z=\sum_{1 \leq i<j \leq p}\left(E_{i j}^{\alpha}\right)^{\top} X_{i j} E_{i j}^{\alpha} \tag{1}
\end{equation*}
$$

for some $X_{i j} \in \mathbb{S}_{+}^{k_{i}+k_{j}}$ and with $E_{i j}^{\alpha} \in \mathbb{R}^{\left(k_{i}+k_{j}\right) \times n}$ being an index matrix.


Figure: Illustration of block factor-width-two decomposition (1). The (i,j) black square represents a submatrix of dimension $k_{i} \times k_{j}, i, j=1,2,3$.

## A hierarchy of inner approximations

A finer partition (or subpartition) of $\alpha=\left\{k_{1}, k_{2}, \ldots, k_{p}\right\}$ is a partition that breaks some blocks of $\alpha$ into smaller blocks.

- Let $\alpha=\{1,1,1,1,1,1\}, \beta=\{2,2,2\}$ and $\gamma=\{4,2\}$. Denote $\alpha \sqsubseteq \beta \sqsubseteq \gamma$.


## Theorem

Let $\alpha \sqsubseteq \beta \sqsubseteq \gamma$ be partitions of $n$ with $\gamma=\left\{\gamma_{1}, \gamma_{2}\right\}$, and let $\mathbf{1}=\{1, \ldots, 1\}$ denote the uniform unit partition. Then,

$$
\mathcal{S D \mathcal { D } _ { n }}=\mathcal{F} \mathcal{W}_{1,2}^{n} \subseteq \mathcal{F} \mathcal{W}_{\alpha, 2}^{n} \subseteq \mathcal{F} \mathcal{W}_{\beta, 2}^{n} \subseteq \mathcal{F} \mathcal{W}_{\gamma, 2}^{n} \equiv \mathbb{S}_{+}^{n}
$$

This flexibility of $\mathcal{F} \mathcal{W}_{\alpha, 2}^{n}$ improves the two drawbacks of $\mathcal{S D} \mathcal{D}_{n}$. As the number $p$ in a partition $\alpha$ decreases:

- First - the approximation quality improves: the largest distance between a unit-norm matrix in $\mathbb{S}_{+}^{n}$ and the cone $\mathcal{F} \mathcal{W}_{\alpha, 2}^{n}$ satisfies

$$
\operatorname{dist}\left(\mathbb{S}_{+}^{n}, \mathcal{F} \mathcal{W}_{\alpha, 2}^{n}\right) \leq \frac{p-2}{p}
$$

- Second - the number of blocks in the summation decreases $\binom{p}{2}$ vs. $\binom{n}{2}$


## Example

Consider the $5 \times 5$ matrix

$$
P(x, y)=\left[\begin{array}{ccccc}
1+6 x+4 y & 3 x+y & 2 x+y & x+4 y & 3 x+3 y \\
3 x+y & 1+6 y & 5 x+3 y & y & 2 x+2 y \\
2 x+y & 5 x+3 y & 1+2 x+2 y & x+2 y & 5 x+6 y \\
x+4 y & y & x+2 y & 1+2 x & 3 x+3 y \\
3 x+3 y & 2 x+2 y & 5 x+6 y & 3 x+3 y & 1+6 x+2 y
\end{array}\right]
$$

and partitions $1=\{1,1,1,1,1\}, \alpha=\{2,1,1,1\}, \beta=\{2,1,2\}$ and $\gamma=\{2,3\}$.



Figure: Regions of the $(x, y)$ plane for which $P(x, y)$ belongs to the block factor-width-two cones $\mathcal{S D}^{2} \mathcal{D}_{5} \equiv \mathcal{F} \mathcal{W}_{1,2}^{5} \subseteq \mathcal{F} \mathcal{W}_{\alpha, 2}^{5} \subseteq \mathcal{F} \mathcal{W}_{\beta, 2}^{5} \subseteq \mathcal{F} \mathcal{W}_{\gamma, 2}^{5} \equiv \mathbb{S}_{+}^{5}$

- The inclusions of the plotted regions reflect the inclusions of the cones and the order relation $\mathbf{1} \sqsubseteq \alpha \sqsubseteq \beta \sqsubseteq \gamma$.


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## Applications to SDPs

- The cones $\mathcal{F} \mathcal{W}_{\alpha, 2}^{n}$ and its dual $\left(\mathcal{F} \mathcal{W}_{\alpha, 2}^{n}\right)^{*}$ approximate the positive semidefinite cone $\mathbb{S}_{+}^{n}$ from the inside and from the outside
- The approximation improves as the partition $\alpha$ is coarsened.
- This leads to convergent sequences of upper and lower bounds on the optimal value $p^{*}$


## Inner approximations

$$
\begin{aligned}
U_{\alpha}:=\min _{X} & \langle C, X\rangle \\
\text { subject to } & \left\langle A_{i}, X\right\rangle=b_{i}, i=1, \ldots, m \\
& X \in \mathcal{F} \mathcal{W}_{\alpha, 2}^{n}
\end{aligned}
$$

## Outer approximations

$$
\begin{aligned}
L_{\alpha}:=\min _{X} & \langle C, X\rangle \\
\text { subject to } & \left\langle A_{i}, X\right\rangle=b_{i}, i=1, \ldots, m \\
& X \in\left(\mathcal{F} \mathcal{W}_{\alpha, 2}^{n}\right)^{*}
\end{aligned}
$$

## Corollary

Let $\alpha_{1} \sqsubseteq \alpha_{2} \sqsubseteq \ldots \sqsubseteq \alpha_{k}=\left\{\alpha_{k 1}, \alpha_{k 2}\right\}$ be a sequence of partitions of $n$. Then,

$$
L_{\alpha_{1}} \leq \cdots \leq L_{\alpha_{k}}=p^{*}=U_{\alpha_{k}} \leq \cdots \leq U_{\alpha_{1}}
$$

## Applications to SOS optimization

- Sum-of-squares (SOS) polynomials: $p(x)$ can be represented as

$$
p(x)=\sum_{i=1}^{m} f_{i}^{2}(x)
$$

- SDP characterization (Parrilo, Lasserre etc.): $p(x)$ is SOS if and only if

$$
p(x)=v_{d}(x)^{T} Q v_{d}(x), \quad Q \succeq 0
$$

where $v_{d}(x)$ is the standard monomial basis.

- SOS: $p(x)=v_{d}(x)^{T} Q v_{d}(x): Q$ is PSD $\longrightarrow$ SDP
- SDSOS: $p(x)=v_{d}(x)^{T} Q v_{d}(x): Q$ is sdd $\longrightarrow$ SOCP
- DSOS: $p(x)=v_{d}(x)^{T} Q v_{d}(x): Q$ is dd $\longrightarrow$ LP
- $\alpha$-SDSOS: $p(x)=v_{d}(x)^{T} Q v_{d}(x): Q$ is $\mathcal{F} \mathcal{W}_{\alpha, 2}^{N} \quad \longrightarrow$ SDP with small blocks


## Numerical examples

Consider a scalar polynomial optimization problem:

$$
\begin{aligned}
\min _{\gamma} & \gamma \\
\text { subject to } & p(x)+\gamma \geq 0, \forall x \in \mathbb{R}^{n}
\end{aligned}
$$

Table: Computational results using SOS and $\alpha$-SDSOS relaxations.

| Full | Number of blocks $p$ in partition $\alpha$ |  |  |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | :---: | :---: | :---: | :---: |
|  | SDP | 10 |  |  |  |  |  | 20 | 50 |
|  | Computational time (seconds) |  |  |  |  |  |  |  |  |  |
|  | 27.3 | 23.3 | 15.6 | 10.1 | 5.36 |  |  |  |  |
|  | 489 | 252 | 98.1 | 66.8 | 28.1 |  |  |  |  |
|  | $\infty$ | 1970 | 783 | 571 | 132 |  |  |  |  |
| 30 | $\infty$ | $\infty$ | 5680 | 3710 | 840 |  |  |  |  |
| Objective values $\gamma$ |  |  |  |  |  |  |  |  |  |
| 15 | -0.92 | -0.75 | 80.1 | 459 | 2240 |  |  |  |  |
| 20 | -0.87 | -0.87 | -0.11 | 251 | 1910 |  |  |  |  |
| 25 | $\infty$ | -1.07 | -0.21 | 231 | 1360 |  |  |  |  |
| 30 | $\infty$ | $\infty$ | -0.37 | 177 | 1770 |  |  |  |  |

## Numerical results

Consider a matrix SOS program

$$
\begin{aligned}
\min _{\gamma} & \gamma \\
\text { subject to } & P(x)+\gamma I \succeq 0, \quad \forall x \in \mathbb{R}^{3},
\end{aligned}
$$

where $P(x)$ is an $r \times r$ polynomial matrix with each element being a quartic polynomial in three variables.

Table: Computational results using SOS, $\alpha$-SDSOS and SDSOS relaxations.

| $r$ | 25 | 30 | 35 | 40 | 45 | 50 |
| :--- | :---: | :---: | ---: | ---: | ---: | ---: |
| Computational time | (seconds) |  |  |  |  |  |
| SOS | 14.4 | 35.9 | 87.2 | 175.0 | 316.0 | 487.8 |
| $\alpha$-SDSOS | 10.8 | 16.6 | 25.3 | 36.0 | 57.4 | 71.4 |
| SDSOS | 1.1 | 1.3 | 1.6 | 2.1 | 2.6 | 3.3 |
| Objective value $\gamma$ |  |  |  |  |  |  |
| SOS | 266.5 | 316.2 | 460.8 | 562.0 | 746.9 | 919.8 |
| $\alpha$-SDSOS | 266.5 | 316.2 | 460.8 | 562.0 | 746.9 | 919.8 |
| SDSOS | 270.3 | 324.8 | 477.7 | 570.9 | 762.2 | 961.7 |

## An iterative scheme

- Ahmadi and Hall ${ }^{4}$ introduced an iterative method over $\mathcal{S D} \mathcal{D}_{n}$ and $\mathcal{D} \mathcal{D}_{n}$.
- This iterative scheme can be naturally extended to the class of block factor-width-two matrices.


Figure: Feasible regions of inner approximations using $\mathcal{D D _ { n }}, \mathcal{S D} \mathcal{D}_{n}$, and

$$
\mathcal{F} \mathcal{W}_{\alpha, 2}^{n}
$$

[^2]
## An iterative scheme



Figure: Inner/Outer approximations of the SDP (2) using $\mathcal{D D}_{n}, \mathcal{S D D}_{n}$, and $\mathcal{F} \mathcal{W}_{2, \alpha}^{n}$ with $\alpha=\{2, \ldots, 2\}^{5}$.

[^3]
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> Introduction: inner/outer approximations for SDPs

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## Summary

A new class of block factor-width-two matrices


Figure: Illustration of (block) factor-width-two decomposition

- It offers a new hierarchy of inner/outer approximations of PSD matrices

$$
\mathcal{S D D}_{n}=\mathcal{F} \mathcal{W}_{1,2}^{n} \subseteq \mathcal{F} \mathcal{W}_{\alpha, 2}^{n} \subseteq \mathcal{F} \mathcal{W}_{\beta, 2}^{n} \subseteq \mathcal{F} \mathcal{W}_{\gamma, 2}^{n} \equiv \mathbb{S}_{+}^{n}
$$

- Applications to inner/outer approximations of SDPs and SOS optimization


## Future work

- Develop special algorithms (such as non-symmetric IPM, non-smooth optimization algorithms) for problems with factor-width-two matrices.


## Thank you for your attention!

## Q \& A

- Yang Zheng, Aivar Sootla, and Antonis Papachristodoulou. "Block factor-width-two matrices and their applications to semidefinite and sum-of-squares optimization." IEEE Transactions on Automatic Control (2022).
- Feng-Yi Liao, and Yang Zheng. "Iterative Inner/outer Approximations for Scalable Semidefinite Programs using Block Factor-width-two Matrices." arXiv preprint arXiv:2204.06759 (2022).


[^0]:    ${ }^{1}$ Zheng, Fantuzzi, and Papachristodoulou. "Chordal and factor-width decompositions for scalable semidefinite and polynomial optimization." Annual Reviews in Control 52 (2021): 243-279.
    ${ }^{2}$ Majumdar, Hall, and Ahmadi. "Recent scalability improvements for semidefinite programming with applications in machine learning, control, and robotics." Annual Review of Control, Robotics, and Autonomous Systems 3 (2020): 331-360.

[^1]:    ${ }^{3}$ Ahmadi, and Majumdar. "DSOS and SDSOS optimization: more tractable alternatives to sum of squares and semidefinite optimization." SIAM J. Appl. Algebra Geom 3.2 (2019): 193-230.

[^2]:    ${ }^{4}$ Ahmadi and Hall. "Sum of squares basis pursuit with linear and second order cone programming." Algebraic and geometric methods in discrete mathematics 2017: 27-53.

[^3]:    ${ }^{5}$ Liao, Feng-Yi, and Yang Zheng. "Iterative Inner/outer Approximations for Scalable Semidefinite Programs using Block Factor-width-two Matrices." arXiv preprint arXiv:2204.06759 (2022).

